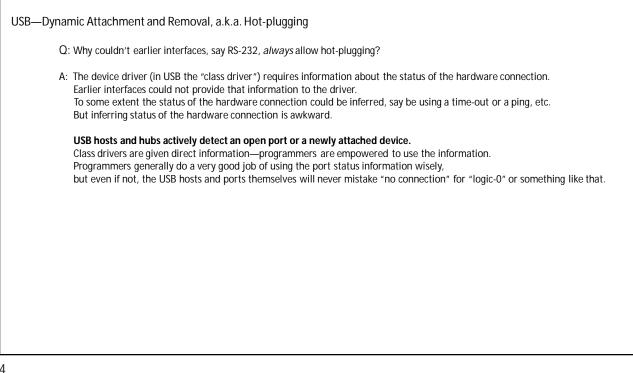
USB—Packet oriented, time-division multiplexed under the host's control. USB transactions each consist of a sequence of packet transmissions under state-machine control. Token packet: From host to everyone (Addressing information, read or write—data payload direction) Payload: as prescribed by token (The data to be communicated. This packet is optional!) Status packet: Receiving device acknowledges error-free reception. This packet is not always present. Each of the three packet types have a sequence of bits, much like an Ethernet frame Sync: opportunity for receiver clock to lock onto the transmitter's clock Packet ID: Defines the protocol for the remaining packets Data: the actual token, payload, or status information Address and endpoint: destination of the packet, distribution of the packet within the destination CRC: error detections (retry the whole transaction if there was an error) EOP: flags the end of the packet

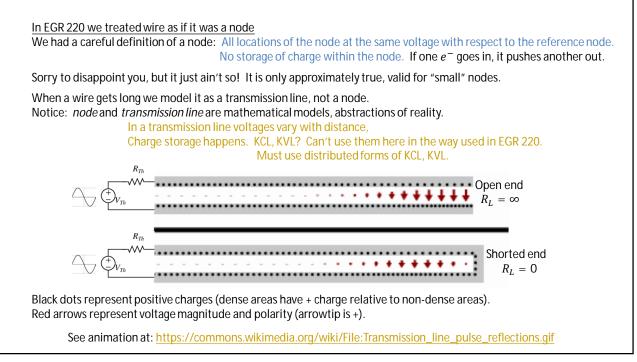
A trans	action is a sequence of packets sent for a particular purpose
Control	transaction
	Mainly for device enumeration
	May also be used for such status matters as, "printer out of paper."
Interrup	ot transaction
	Not what it seems. In USB there is no way for a device to grab the attention of the host 😕 Scheduled polling—host periodically polls the device
	Typically used for low-speed HID devices: A character was typed, or the mouse moved, etc.
	If the device has an interrupt request, that gets communicated when the device is polled.
	There is a bound on the latency. (But it is nowhere as short as a hardware interrupt can produce.)
Isochroi	nous transaction
	Intended for media transmission—smooth motion and sound are needed
	Time-slotted packets provide guaranteed bandwidth (bit-rate)
	Bounded latency—keeps up with "real time."
	Error detection. If an error, discard the packet. No retry. Receiver must "make up" something. Not available in low speed.
Bulk tra	nsaction
Duntitu	Intended for file transfer: A burst of data as fast as possible, then long intervals of silence.
	Error detection with retry as much as needed for 100% valid transfer of data
	packets have lowest priority relative to above types. No guarantee of bandwidth or latency. Unidirectional, not available in low speed.

USB—Dynamic Attachment and Removal, a.k.a. Hot-plugging		
We are so used to it that this hardly seems to be a "feature." When you insert a USB plug (anywhere) or pull one out something sensible happens!		
Things do not "crash" when you pull a USB plug—only the device associated with the plug stops working. Believe it or not, this was a first with USB.		
Interfaces of the 1990's era and earlier were subject to capricious behavior if the hardware connection was lost.		
Typical bad behavior of earlier data communications:		
E.g. if you unplug the printer the printer driver starts filling up memory with pages to print. Eventually the system's memory fills up		
Eventually the operating system can no longer work.		
The whole system, every user, every program in progress, crashes.		
What we want:		
The printer's USB cable gets unplugged for some reason. The printer driver starts fills up, but it is given only a limited segment of memory		
This printer driver recognizes its memory is full and signals all programs that write to it—offline.		
Problem is limited, other than printer services, nothing else is impacted. This "good" behavior is designed into the "class drivers."		
This good behavior is required by the USB standard.		
Still, hot-plugging is not recommended. E.g. you are writing to a USB thumb drive while you unplug it. This could leave behind corrupted data. I rare cases it could even brick the thumb drive. (Brick: Cause the lost of internal programming so that it will for all programming so that it will for all programming so that it will for all programming so that it will be the "sofely remove USP douber" ("unmount" in Univ. Journal of your OS		
practical purposes never be useful again.) Use the "safely remove USB device" ("unmount" in Unix, Linux) option of your OS.		

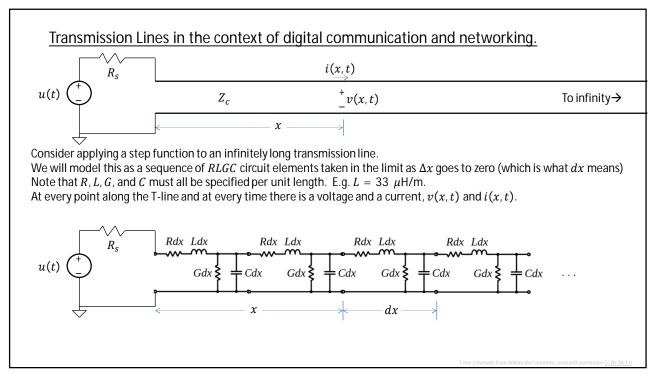


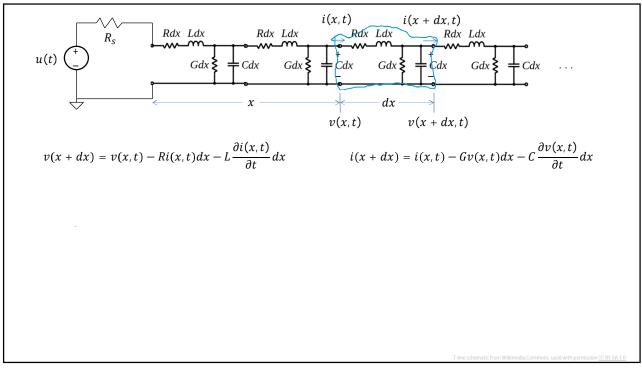


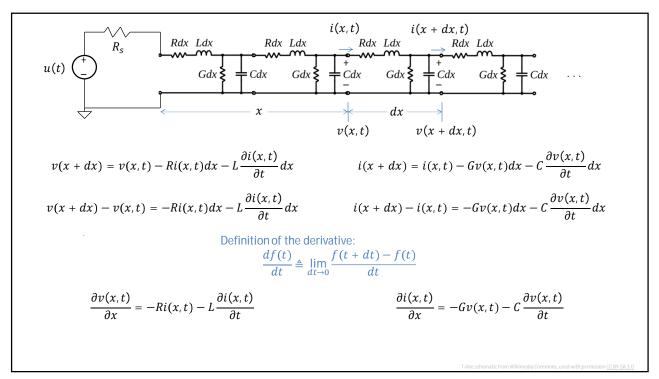


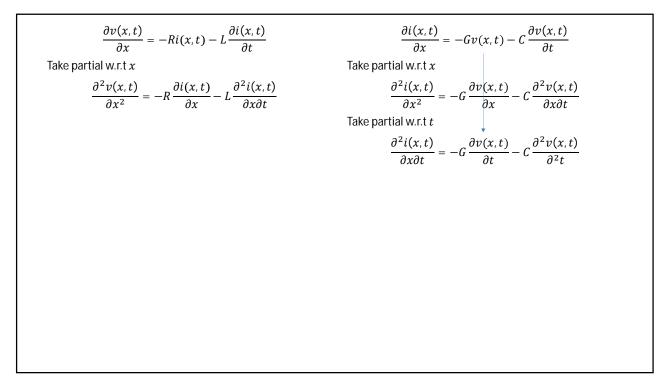


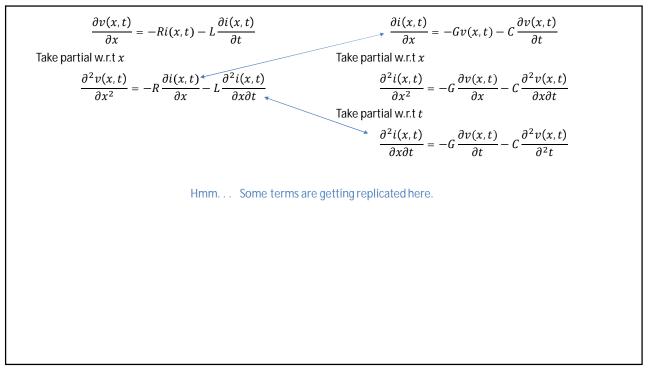


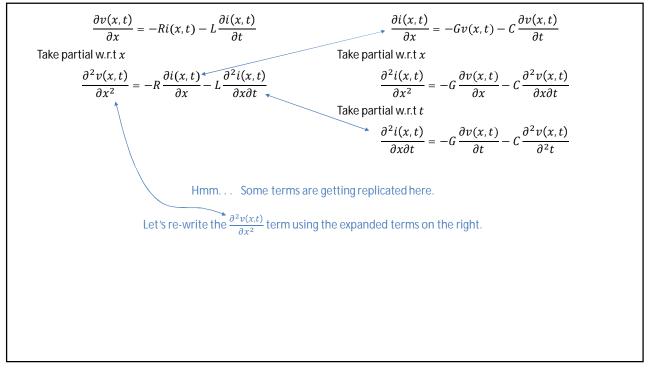


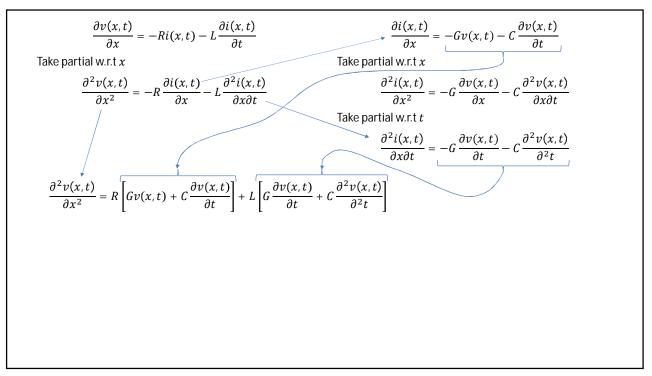


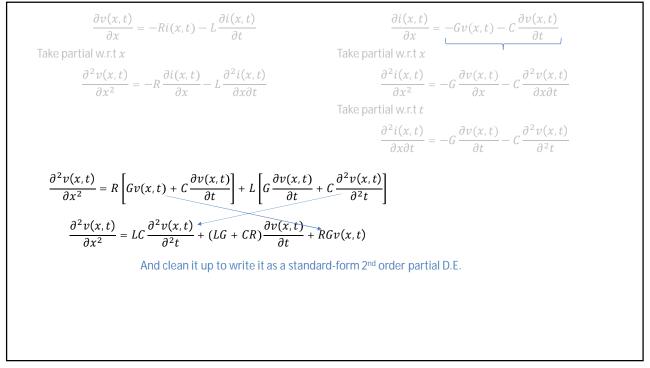


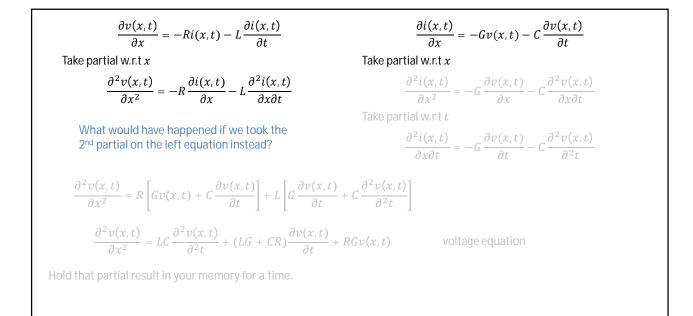


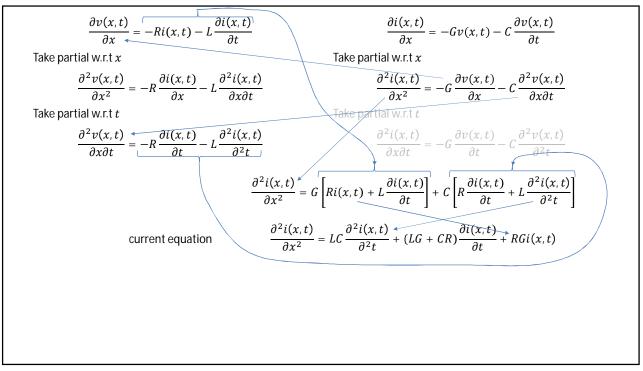












The two equations we have derived are called the Telegrapher's Equations

$$\frac{\partial^2 v(x,t)}{\partial x^2} = LC \frac{\partial^2 v(x,t)}{\partial^2 t} + (LG + CR) \frac{\partial v(x,t)}{\partial t} + RGv(x,t)$$

$$\frac{\partial^2 i(x,t)}{\partial x^2} = LC \frac{\partial^2 i(x,t)}{\partial^2 t} + (LG + CR) \frac{\partial i(x,t)}{\partial t} + RGi(x,t)$$
For short transmission lines such as are usually encountered in microcontroller situations we can assume that the transmission line is lossless. That is, $R = 0$ and $G = 0$.

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Wave Equations

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Let $v(x, t) = \cos(\omega(t - x/V_P))$ where ω = frequency (rad/sec) and V_P = propagation speed (m/s)